

# VALLIAMMAI ENGINEERING COLLEGE

SRM Nagar, Kattankulathur – 603 203.

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING



## QUESTION BANK

**SUBJECT : CS8501**

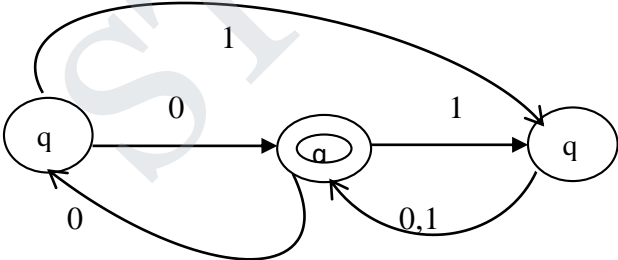
**SEM / YEAR: V/III**

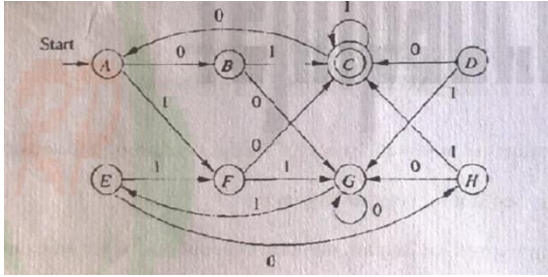
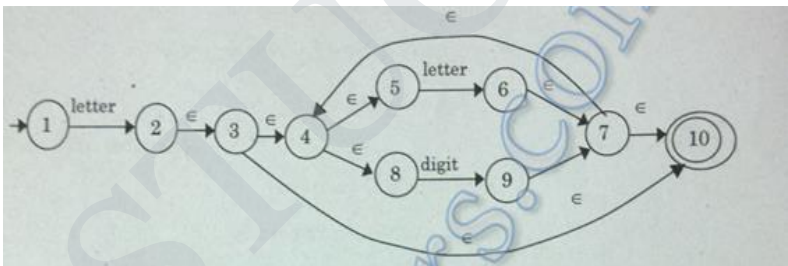
UNIT I		AUTOMATA FUNDAMENTALS	
Introduction to formal proof – Additional forms of Proof – Inductive Proofs – Finite Automata – Deterministic Finite Automata – Non-deterministic Finite Automata – Finite Automata with Epsilon Transitions			
PART – A			
Q.No	Questions	BT Level	Competence
1.	<b>Differentiate</b> between DFA and NFA.	BTL-2	Understand
2.	<b>Define</b> DFA	BTL-1	Remember
3.	<b>Define</b> inductive proof.	BTL-1	Remember
4.	<b>Identify</b> NFA- $\epsilon$ to represent $a^*b   c$	BTL-1	Remember
5.	<b>Consider</b> the String $X=110$ and $y=0110$ find i) $XY$ ii) $X^2$ iii) $YX$ iv) $Y^2$	BTL-4	Analyze
6.	<b>Describe</b> the following language over the input set $A=\{a,b\}$ $L=\{ a^n b^n   n \geq 1 \}$	BTL-4	Analyze
7.	<b>Describe</b> what is non-deterministic finite automata and the applications of automata theory.	BTL-1	Remember
8.	Illustrate the induction principle.?	BTL-3	Apply
9.	What is proof by contradiction ?	BTL-1	Remember
10.	<b>Describe</b> an identifier with a transition diagram (automata).	BTL-2	Understand
11.	<b>Define</b> $\epsilon$ -NFA	BTL-1	Remember
12.	<b>Summarize</b> minimization of DFA	BTL-5	Evaluate
13.	<b>Give</b> the non-deterministic automata to accept strings containing the substring 0101	BTL-2	Understand
14.	<b>Illustrate</b> if $L$ be a set accepted by an NFA then there exists a DFA that accepts $L$ .	BTL-3	Apply
15.	<b>Define</b> the term epsilon transition.	BTL-2	Understand
16.	<b>Summarize</b> the extended transition function for a $\epsilon$ -NFA	BTL-5	Evaluate

17.	<b>Create</b> a FA which accepts the only input 101 over the input set $Z=\{0,1\}$	BTL-6	Create
18.	Describe a Finite automata and give its types.	BTL-4	Analyze
19.	<b>Illustrate</b> deductive proof.	BTL-3	Apply
20.	<b>Create</b> a FA which checks whether the given binary number is even.	BTL-6	Create

**PART - B**

1.	<p>(i)<b>Explain</b> if L is accepted by an NFA with <math>\epsilon</math>-transition then show that L is accepted by an NFA without <math>\epsilon</math>-transition.(6)</p> <p>(ii)<b>Construct</b> a DFA equivalent to the NFA. <math>M=(\{p,q,r\},\{0,1\},\delta,p,\{q,s\})</math> Where <math>\delta</math> is defined in the following table.(7)</p> <table border="1" style="margin-left: auto; margin-right: auto;"> <tr> <td></td> <td>0</td> <td>1</td> </tr> <tr> <td>p</td> <td>{q,s}</td> <td>{q}</td> </tr> <tr> <td>q</td> <td>{r}</td> <td>{q,r}</td> </tr> <tr> <td>r</td> <td>{s}</td> <td>{p}</td> </tr> <tr> <td>s</td> <td>-</td> <td>{p}</td> </tr> </table>		0	1	p	{q,s}	{q}	q	{r}	{q,r}	r	{s}	{p}	s	-	{p}	BTL-5	Evaluate
	0	1																
p	{q,s}	{q}																
q	{r}	{q,r}																
r	{s}	{p}																
s	-	{p}																
2.	<b>Prove</b> for every $n \geq 1$ by mathematical induction $\sum (i)=\{n(n+1)/2\}$ . (13)	BTL-3	Apply															
3.	Let L be a set accepted by a NFA then <b>show</b> that there exists a DFA that accepts L.(13)	BTL-1	Remember															
4.	<b>Give</b> the NFA that accepts all strings that end in 01. Give its transition table and the extended transition function for the input string 00101. Also construct a DFA for the above NFA using subset construction method.(13)	BTL-2	Understand															
5.	<p>Construct DFA equivalent to the NFA given below: (13)</p>	BTL-2	Understand															

<p>6.</p>	<p>(i)<b>Compose</b> that a language L is accepted by some <math>\epsilon</math>-NFA if and only if L is accepted by some DFA. (6)</p> <p>(ii)<b>Consider</b> the following <math>\epsilon</math>-NFA. Compute the <math>\epsilon</math>-closure of each state and find its equivalent DFA. (7)</p> <table border="1" data-bbox="306 293 794 490"> <tr> <td></td> <td><math>\epsilon</math></td> <td>a</td> <td>b</td> <td>C</td> </tr> <tr> <td><math>\rightarrow p</math></td> <td><math>\phi</math></td> <td>{p}</td> <td>{q}</td> <td>{r}</td> </tr> <tr> <td>q</td> <td>{p}</td> <td>{q}</td> <td>{r}</td> <td><math>\phi</math></td> </tr> <tr> <td>*r</td> <td>{q}</td> <td>{r}</td> <td><math>\phi</math></td> <td>{p}</td> </tr> </table>		$\epsilon$	a	b	C	$\rightarrow p$	$\phi$	{p}	{q}	{r}	q	{p}	{q}	{r}	$\phi$	*r	{q}	{r}	$\phi$	{p}	<p>BTL-6</p>	<p>Create</p>
	$\epsilon$	a	b	C																			
$\rightarrow p$	$\phi$	{p}	{q}	{r}																			
q	{p}	{q}	{r}	$\phi$																			
*r	{q}	{r}	$\phi$	{p}																			
<p>7.</p>	<p>i)<b>Classify</b> how a language L is accepted by some DFA if L is accepted by some NFA(7)</p> <p>(ii)Convert the following NFA to its equivalent DFA.(6)</p> <table border="1" data-bbox="360 656 896 869"> <tr> <td></td> <td>0</td> <td>1</td> </tr> <tr> <td>p</td> <td>{p,q}</td> <td>{p}</td> </tr> <tr> <td>q</td> <td>{r}</td> <td>{r}</td> </tr> <tr> <td>r</td> <td>{s}</td> <td><math>\phi</math></td> </tr> <tr> <td>*s</td> <td>{s}</td> <td>{s}</td> </tr> </table>		0	1	p	{p,q}	{p}	q	{r}	{r}	r	{s}	$\phi$	*s	{s}	{s}	<p>BTL-3</p>	<p>Apply</p>					
	0	1																					
p	{p,q}	{p}																					
q	{r}	{r}																					
r	{s}	$\phi$																					
*s	{s}	{s}																					
<p>8.</p>	<p>i)<b>Construct</b> the DFA to recognize odd number of 1's and even number 0's (7)</p> <p>ii) Construct the DFA over {a,b} which produces not more than 3 a's (6)</p>	<p>BTL-1</p>	<p>Remember</p>																				
<p>9.</p>	<p>(i)<b>Point out</b> the steps in conversion of NFA to DFA and for the following convert NFA to a DFA(7)</p> <table border="1" data-bbox="217 1211 472 1391"> <tr> <td><math>\delta</math></td> <td>a</td> <td>b</td> </tr> <tr> <td>p</td> <td>{p}</td> <td>{p,q}</td> </tr> <tr> <td>q</td> <td>{r}</td> <td>{r}</td> </tr> <tr> <td>r</td> <td>{<math>\phi</math>}</td> <td>{<math>\phi</math>}</td> </tr> </table> <p>(ii)<b>Infer</b> the following to a regular expression.(6)</p> 	$\delta$	a	b	p	{p}	{p,q}	q	{r}	{r}	r	{ $\phi$ }	{ $\phi$ }	<p>BTL-4</p>	<p>Analyze</p>								
$\delta$	a	b																					
p	{p}	{p,q}																					
q	{r}	{r}																					
r	{ $\phi$ }	{ $\phi$ }																					

<p>10.</p>	<p><b>Identify</b> and explain the algorithm for minimization of DFA. Using the above algorithm minimize the following DFA.(13)</p> 	<p>BTL-1</p>	<p>Remember</p>																				
<p>11.</p>	<p><b>Tabulate</b> the difference between the NFA and DFA .Convert the following <math>\epsilon</math>-NFA to DFA.(13)</p> <table border="1" data-bbox="411 779 842 1008"> <thead> <tr> <th>states</th> <th><math>\epsilon</math></th> <th>a</th> <th>b</th> <th>c</th> </tr> </thead> <tbody> <tr> <td>p</td> <td><math>\Phi</math></td> <td>{p}</td> <td>{q}</td> <td>{r}</td> </tr> <tr> <td>q</td> <td>{p}</td> <td>{q}</td> <td>{r}</td> <td><math>\Phi</math></td> </tr> <tr> <td>*r</td> <td>{q}</td> <td>{r}</td> <td><math>\phi</math></td> <td>{p}</td> </tr> </tbody> </table>	states	$\epsilon$	a	b	c	p	$\Phi$	{p}	{q}	{r}	q	{p}	{q}	{r}	$\Phi$	*r	{q}	{r}	$\phi$	{p}	<p>BTL-1</p>	<p>Remember</p>
states	$\epsilon$	a	b	c																			
p	$\Phi$	{p}	{q}	{r}																			
q	{p}	{q}	{r}	$\Phi$																			
*r	{q}	{r}	$\phi$	{p}																			
<p>12.</p>	<p>(i).<b>Describe</b> the extended transition function for NFA ,DFA and <math>\epsilon</math>-NFA(6)</p> <p>(ii) Consider the following <math>\epsilon</math>-NFA for an identifier .Consider the <math>\epsilon</math>-closure of each state and <b>give</b> it's equivalent DFA.(7)</p> 	<p>BTL-2</p>	<p>Understand</p>																				
<p>13.</p>	<p>(i)<b>Given</b> <math>\Sigma = \{a,b\}</math> Analyze and construct a DFA which recognize the language <math>L = \{b^m a b^n : m,n &gt; 0\}</math> (13)</p>	<p>BTL-4</p>	<p>Analyze</p>																				
<p>14.</p>	<p>(i)<b>Analyze</b> and Prove that if n is a positive integer such that n mod 4 is 2 or 3 then n is not a perfect square.(6)</p> <p>(ii)<b>Construct</b> a DFA that accept the string {0,1} that always ends with 00 (7)</p>	<p>BTL-4</p>	<p>Analyze</p>																				

PART – C

1.	<p>(i) <b>Draw</b> and Explain the transition diagram for recognizing the set of all operators in c Language.(8)</p> <p>(ii)<b>Evaluate</b> a DFA from the given NFA(7)  <math>M=(\{q_0,q_1\},\{a,b\},\delta,q_0,\{q_1\})</math> with the state table diagram for <math>\delta</math> given below:</p> <table border="1" data-bbox="156 353 699 510"> <tr> <td><math>\delta</math></td> <td>a</td> <td>b</td> </tr> <tr> <td><math>q_0</math></td> <td><math>\{q_0,q_1\}</math></td> <td><math>q_1</math></td> </tr> <tr> <td><math>q_1</math></td> <td><math>\Phi</math></td> <td><math>\{q_0,q_1\}</math></td> </tr> </table>	$\delta$	a	b	$q_0$	$\{q_0,q_1\}$	$q_1$	$q_1$	$\Phi$	$\{q_0,q_1\}$	BTL-5	Evaluation											
$\delta$	a	b																					
$q_0$	$\{q_0,q_1\}$	$q_1$																					
$q_1$	$\Phi$	$\{q_0,q_1\}$																					
2.	<p>Construct the following <math>\epsilon</math>-NFA to DFA.(15)</p> <table border="1" data-bbox="363 586 842 810"> <tr> <td>states</td> <td><math>\epsilon</math></td> <td>a</td> <td>b</td> <td>c</td> </tr> <tr> <td>p</td> <td><math>\Phi</math></td> <td><math>\{p\}</math></td> <td><math>\{q\}</math></td> <td><math>\{r\}</math></td> </tr> <tr> <td>q</td> <td><math>\{p\}</math></td> <td><math>\{q\}</math></td> <td><math>\{r\}</math></td> <td><math>\{p,q\}</math></td> </tr> <tr> <td>*r</td> <td><math>\{q\}</math></td> <td><math>\{r\}</math></td> <td><math>\phi</math></td> <td><math>\Phi</math></td> </tr> </table>	states	$\epsilon$	a	b	c	p	$\Phi$	$\{p\}$	$\{q\}$	$\{r\}$	q	$\{p\}$	$\{q\}$	$\{r\}$	$\{p,q\}$	*r	$\{q\}$	$\{r\}$	$\phi$	$\Phi$	BTL-6	Create
states	$\epsilon$	a	b	c																			
p	$\Phi$	$\{p\}$	$\{q\}$	$\{r\}$																			
q	$\{p\}$	$\{q\}$	$\{r\}$	$\{p,q\}$																			
*r	$\{q\}$	$\{r\}$	$\phi$	$\Phi$																			
3.	<p><b>Infer</b> the DFA which is accepting the following language over the alphabet<math>\{0,1\}</math>.The set of all the strings beginning with a1 that when interrupted as a binary integer , is multiple of 5, For example strings 101,1010 and 1111 are in the language 0,100 and 111 are not.(15)</p>	BTL-4	Analyze																				
4.	<p><b>Rewrite</b> the basic approach to convert from NFA to regular expression. Illustrate with an example(15)</p>	BTL-6	Create																				

**UNIT II REGULAR EXPRESSION AND LANGUAGES**

Regular Expressions – FA and Regular Expressions – Proving Languages not to be regular – Closure Properties of Regular Languages – Equivalence and Minimization of Automata.

**PART - A**

Q.No	Questions	BT Level	Competence
1.	<b>List</b> the operators of Regular Expressions	BTL-1	Remember
2.	<b>Differentiate</b> between regular expression and regular	BTL-1	Remember
3.	<b>Tabulate</b> the regular expression for the following L1=set of strings 0 and 1 ending in 00	BTL-4	Analyze
4.	<b>What</b> are the closure properties of regular languages?	BTL-2	Understand
5.	<b>Explain</b> a finite automaton for the regular expression $0^*1^*$ .	BTL-1	Remember
6.	<b>Illustrate</b> a regular expression for the set of all the strings	BTL-1	Remember

7.	<b>Illustrate</b> a regular expression for the set of all the strings have odd number of 1's R.E= $1(0+1)^*$	BTL-3	Apply
8.	<b>Compose</b> the difference between the + closure and * closure	BTL-4	Analyze
9.	<b>Illustrate</b> a regular expression for the set of all strings of 0's	BTL-2	Understand
10.	<b>What</b> is the Closure property of regular set S.?	BTL-2	Understand
11.		BTL-2	Understand
12.	<b>Find</b> out the language generated by the regular expression $(0+1)^*$ .	BTL-5	Evaluate
13.	<b>Name</b> the four closure properties of RE.	BTL-1	Remember
14.	Is it true the language accepted by any NFA is different from the regular language? <b>Justify</b> your answer.	BTL-4	Analyze
15.	<b>Show</b> the complement of a regular language is also regular.	BTL-3	Apply
16.	<b>Construct</b> a DFA for the regular expression $aa^*bb^*$ .	BTL-3	Apply
17.	<b>State</b> the precedence of RE operator.	BTL-5	Evaluate
18.	<b>Construct</b> RE for the language over the set $z=\{a,b\}$ in which total number of a's are divisible by 3.	BTL-6	Create
19.	Define RE.	BTL-1	Remember
20.	<b>Create</b> RE to describe an identifier and positive integer.	BTL-6	Create
<b>PART - B</b>			
1.	<b>Demonstrate</b> how the set $L = \{ab^n/n \geq 1\}$ is not a regular.(13)	BTL5	Evaluate
2.	<b>Express</b> that the regular languages are closed under:(13) (a)union (b)intersection(c)Kleene Closure(d)Complement (e)Difference	BTL-1	Remember
3.	<b>Examine</b> whether the language $L=(0^n1^n   n \geq 1)$ is regular or not? Justify your answer (13)	BTL-2	Understand
4.	(i) <b>Describe</b> a Regular Expression. Write a Regular Expression for the set of strings that consists of alternating 0's and 1's.(6)  (ii) <b>Construct</b> Finite Automata equivalent to the regular expression $(ab+a)^*(7)$ .	BTL1	Remember
5.	(i) <b>Describe</b> the closure properties of regular languages.(6)  (ii) <b>Describe</b> NFA with epsilon for the RE= $(a/b)^*ab$ and convert it into DFA and further find the minimized DFA.(7)	BTL1	Remember

6.	<b>Demonstrate</b> how the set $L = \{a^n b^n / n \geq 0\}$ is not a regular. (13)	BTL-3	apply
7.	<b>Verify</b> the whether $L = \{a^{2n} / n \geq 1\}$ regular (13)	BTL-3	Apply
8.	i) Prove The reverse of a regular language is regular (6) ii) A homomorphism of regular language is regular (7)	BTL-4	Analyze
9.	Discuss on regular expressions (13)	BTL-2	Understand
10	<b>Construct</b> NDFFA for given RE using Thomson rule. (13) i) $a.(a+b)^* ab$ ii) $(a.b)^*$ iii) $(a+b)$	BTL-6	Create
11	Explain the DFA Minimization algorithm with an example. (13)	BTL-1	Remember
12	<b>Demonstrate</b> how the set $L = \{a^n b^m / m, n \geq 1\}$ is not a regular. (13)	BTL 2	Understand
13	i) <b>Prove</b> the $L_1$ and $L_2$ are two languages then $L_1 - L_2$ is regular (7) ii) <b>Prove</b> the $L_1$ and $L_2$ are two languages then $L_1 \cdot L_2$ is regular (6)	BTL 4	Analyze
14	i) <b>Prove</b> the $L_1$ and $L_2$ are two languages then $L_1 \cup L_2$ is regular (7) ii) Prove the $L_1$ and $L_2$ are two languages then $L_1$ intersection $L_2$ is regular (6)	BTL-4	Analyze

**PART – C**

1.	<p>(i) Deduce into regular expression that denotes the language accepted by following DFA. (7)</p> <pre> graph LR     start(( )) --&gt; 1((1))     1 -- 1 --&gt; 1     1 -- 0 --&gt; 2(((2)))     2 -- "0,1" --&gt; 2     </pre> <p>(ii) <b>Evaluate</b> the equalities for the following RE and prove for the same (8)</p> <ol style="list-style-type: none"> <li><math>b+ab^* +aa^*b+aa^*ab^*</math></li> <li><math>a^*(b+ab^*)</math>.</li> <li><math>a(a+b)^*+aa(a+b)^*+aaa(a+b)^*</math></li> </ol>	BTL-5	Evaluate
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2.	Set the algorithm for minimization of a DFA. <b>Develop</b> a minimized DFA for the RE $(a+b)(a+b)^*$ and trace for the string baaaab.(15)	BTL-6	Create
3.	<b>Point out</b> about the regular expression and regular Language. (15)	BTL-4	Analyze
4.	<b>Develop</b> the procedure for minimizing DFA with example (15)	BTL-6	Create

**UNIT III CONTEXT FREE GRAMMAR AND LANGUAGES**

CFG – Parse Trees – Ambiguity in Grammars and Languages – Definition of the Pushdown Automata – Languages of a Pushdown Automata – Equivalence of Pushdown Automata and CFG, Deterministic Pushdown Automata.

**PART - A**

Q.No	Questions	BT Level	Competence
1.	<b>Express</b> the ways of languages accepted by PDA and define them?	BTL 2	Understand
2.	<b>Summarize</b> PDA .Convert the following CFG to PDA $S \rightarrow aAA, A \rightarrow aS bS a$ .	BTL 2	Understand
3.	<b>Define</b> ambiguous grammar and CFG	BTL 1	Remember
4.	<b>Define</b> parse tree and derivation.	BTL 1	Remember
5.	<b>Examine</b> the context free Grammar representing the set of Palindrome over $(0+1)^*$	BTL 2	Understand
6.	<b>Compare</b> Deterministic and Non deterministic PDA. Is it true that non deterministic PDA is more powerful than that of deterministic PDA? Justify your answer.	BTL 2	Understand
7.	<b>When</b> is PDA said to be deterministic?	BTL 1	Remember
8.	<b>Examine</b> the string aaabbabbba for the Grammar G with $S \rightarrow aB bA$ $A \rightarrow a aS bAA$ $B \rightarrow b bS aBB$	BTL 5	Evaluate
9.	<b>Examine</b> whether a pushdown automata has a memory?	BTL 1	Remember
10.	<b>Design</b> equivalence of PDA and CFG.	BTL 6	Create
11.	<b>Point out</b> the languages generated by a PDA using final state of the PDA and empty stack of that PDA	BTL 4	Analyze
12.	<b>Illustrate</b> the rule for construction of CFG from given PDA	BTL 3	Apply
13.	<b>Give</b> a CFG for the language of palindrome string over $\{a,b\}$ . Write the CFG for the language $L = \{a^n b^n   n \geq 1\}$ .	BTL 5	Evaluate
14.	What is Instantaneous Descriptions ( ID )	BTL 1	Remember
15.	<b>Show</b> that $L = \{a^p   p \text{ is prime}\}$ is not context free.	BTL 3	Apply



16.	<b>Infer</b> the CFG for the set of strings that contains equal number of a's and b's over $\Sigma = \{a,b\}$	BTL 4	Analyze
17.	<b>Point out</b> the various types of grammar with example	BTL 1	Remember
18.	<b>Illustrate</b> the rightmost derivation $(a+b)^*c$ for using the grammar and also state whether a given grammar is ambiguous one or not. $E \rightarrow E+E/E^*E/(E)/id$	BTL 3	Apply
19.	<b>Point out</b> the additional features a PDA has when compared with NFA.	BTL 4	Analyze
20.	<b>Convince</b> your answer of a context free grammar for the given expression $(a+b)(a+b+0+1)^*$	BTL6	Create
<b>PART - B</b>			
1.	(i) <b>Discuss</b> about PDA and CFL Prove the equivalence of PDA and CFL.(6) (ii)If L is Context free language then <b>prove</b> that there exists PDA M such that $L=N(M)$ . (7)	BTL 2	Understand
2.	(i) <b>Describe</b> the different types of acceptance of a PDA. Are they equivalent in sense of language acceptance? Justify your answer. (7) (ii)Design a PDA to accept $\{0^n1^n   n > 1\}$ . Draw the transition diagram for the PDA and <b>identify</b> the instantaneous description(ID)of the PDA which accepts the string '0011.(6)	BTL 1	Remember
3.	(i) <b>Identify that</b> deterministic PDA is less powerful than non nondeterministic PDA.(7) (ii)Construct a PDA accepting $\{a^nb^ma^n / m, n \geq 1\}$ by empty stack. Also <b>tell</b> the corresponding context-free grammar accepting the same set.(6)	BTL 1	Remember
4.	(i) <b>Describe</b> and draw the parse tree for the string $1+2*3$ Given the grammar $G=(V, \Sigma, R, E)$ where $V = \{E, D, 1, 2, 3, 4, 5, 6, 7, 9, 0, +, -, *, /, 9, \}$ $\Sigma = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 0, +, -, *, /, (, )\}$ where R contains the following rules : $E \rightarrow D (E) E+E E-E E/E$ $D \rightarrow 0 1 2 \dots 9$ (6)  (ii)Let $G=(V, T, P, S)$ be a Context Free Grammar then prove that if the recursive inference procedure call tells us that terminal string W is in the language of variable A ,then there is a parse tree with a root A and yield w. (7)	BTL 6	Create

<p>5.</p>	<p><b>(i)Define</b> Non Deterministic Push Down Automata. Is it true that DPDA and NDPDA are equivalent in the sense of language acceptance is concern? Justify Your answer.(5)  <b>(ii)Convert</b> PDA to CFG.PDA is given by  <math>P=(\{p,q\},\{0,1\},\{X,Y\},\delta,q,Z), \delta</math> is defined by  <math>\delta(p,1,z)=\{(p,XZ)\},</math>  <math>\delta(p,\epsilon,Z)=\{(p,\epsilon)\},</math>  <math>\delta(p,1,X)=\{(p,XX)\},</math>  <math>\delta(q,1,X)=\{(q,\epsilon)\},</math>  <math>\delta(p,0,X)=\{(q,X0)\}</math>  <math>\delta(q,0,Z)=\{(p,Z)\}</math> (8)</p>	<p>BTL 1</p>	<p>Remember</p>
<p>6.</p>	<p><b>(i)Define</b> PDA. <b>Give</b> an Example for a language accepted byPDA by empty stack.(7)  <b>(ii)Convert</b> the grammar  <math>S \rightarrow 0S1 A</math>  <math>A \rightarrow 1A0 S  \epsilon</math> into PDA that accepts the same language by the empty stack .Check whether 0101 belongs to <math>N(M)</math>.(6)</p>	<p>BTL 2</p>	<p>Understand</p>
<p>7.</p>	<p><b>(i)Analyze</b> the theorem: If L is Context free language then prove that there exists PDA M such that <math>L=N(M)</math>. (7)  <b>(ii) Prove</b> that if there is PDA that accepts by the final state then there exists an equivalent PDA that accepts by Null State.(6)</p>	<p>BTL 4</p>	<p>Analyze</p>
<p>8.</p>	<p><b>Solve</b> the following grammar  <math>S \rightarrow aAa   bBb   BB</math>  <math>A \rightarrow C</math>  <math>B \rightarrow S   A</math>  <math>C \rightarrow S   \epsilon</math> for the string abaaba .Give  i) Left most derivation(3)  ii)Right most derivation(3)  iii)Derivation Tree(3)  iv) For the string abaabbba, find the right most derivation.(4)</p>	<p>BTL 5</p>	<p>Evaluate</p>
<p>9.</p>	<p><b>(i)Examine</b>Construct the grammar for the following PDAM.  <math>M=(\{q0, q1\},\{0,1\},\{X,z0\},\delta,q0,Z0,\Phi)</math> and where <math>\delta</math> is given by  <math>\delta(q0,0,z0)=\{(q0,XZ0)\},</math>  <math>\delta(q0,0,X)=\{(q0,XX)\},</math>  <math>\delta(q0,1,X)=\{(q1,\epsilon)\},</math>  <math>\delta(q1,1,X)=\{(q1,\epsilon)\},</math>  <math>\delta(q1,\epsilon,X)=\{(q1,\epsilon)\},</math>  <math>\delta(q1,\epsilon,Z0)=\{(q1,\epsilon)\}.</math> (7)  <b>(ii)Prove</b> that if L is <math>N(M1)</math> for some PDA M1 then L is <math>L(M2)</math> for some PDA M2. (6)</p>	<p>BTL 3</p>	<p>Apply</p>
<p>10.</p>	<p>Construct a PDA that recognizes and <b>analyzes</b>the language <math>\{a^i b^j c^k   i,j,k&gt;0 \text{ and } i=j \text{ or } i=k\}.</math>  <b>Explain</b> about PDA acceptance  i) From empty Stack to final state. (6)  ii) From Final state to Empty Stack. (7)</p>	<p>BTL 4</p>	<p>Analyze</p>

<p><b>11.</b></p>	<p><b>Examine</b> and construct a CFG <math>G</math> which accepts <math>N(M)</math>, where <math>M = (\{q_0, q_1\}, \{a, b\}, \{z_0, z\}, \delta, q_0, z_0, \Phi)</math> and where <math>\delta</math> is given by  <math>\delta(q_0, b, z_0) = \{(q_0, zz_0)\}</math>  <math>\delta(q_0, \epsilon, z_0) = \{(q_0, \epsilon)\}</math>  <math>\delta(q_0, b, z) = \{(q_0, zz)\}</math>  <math>\delta(q_0, a, z) = \{(q_1, z)\}</math>  <math>\delta(q_1, b, z) = \{(q_1, \epsilon)\}</math>  <math>\delta(q_1, a, z_0) = \{(q_0, z_0)\}</math>                  Show that <math>a^n b^n c^n</math> is not context free language i.e show that the set of strings of <math>a</math>'s and <math>b</math>'s and <math>c</math>'s with an equal number of each is not context free(13)</p>	<p>BTL-1</p>	<p>Remember</p>
<p><b>12.</b></p>	<p>(i) <b>Describe</b> the PDA that accept the given CFG (7)  <math>S \rightarrow xaax</math>  <math>X \rightarrow ax/bx/\epsilon</math>                  (ii) <b>Express</b> a PDA for the language <math>a^n b^m a^{n+m}</math> (6)</p>	<p>BTL-2</p>	<p>Understand</p>
<p><b>13.</b></p>	<p>(i) Illustrate a PDA for the language <math>\{WCWR/W \in \{0,1\}^*\}</math>. (7)                  (ii) <b>Illustrate</b> a CFG for the constructed PDA. (6)</p>	<p>BTL-3</p>	<p>Apply</p>
<p><b>14.</b></p>	<p>(i) <b>Identify</b> CFG for the language <math>L = \{0^i 1^j 0^k \mid j &gt; i+k\}</math> (7)                  (ii) <b>Define</b> derivation tree. Explain its uses with an example. (6)</p>	<p>BTL-4</p>	<p>Analyze</p>
<p><b>PART - C</b></p>			
<p>1.</p>	<p>(i) Design and <b>Explain</b> a PDA to accept each of the following language  <math>\{a^i b^j c^k \mid i=j \text{ or } j=k\}</math> (7)                  (ii) The set of all string with twice as many 0's and 1's. (8)</p>	<p>BTL-5</p>	<p>Evaluation</p>
<p>2.</p>	<p>(i) Let <math>P</math> be a PDA with empty stack language <math>L = N(P)</math> and suppose that <math>\epsilon</math> is not in <math>L</math>. <b>Design</b> how you would modify <math>P</math> so that it accepts <math>L \cup \{\epsilon\}</math> by empty stack. (8).                  (ii) <b>Design</b> a DPDA for even length palindrome. (7)</p>	<p>BTL-6</p>	<p>Create</p>
<p>3.</p>	<p>(i) Convert the following CFG to PDA and <b>analyze</b> the answer <math>(a+b)</math> and <math>a^{++}</math>. (8)  <math>I \rightarrow a b Ia Ib I0 I1</math>  <math>E \rightarrow I E+E E^*E (E)</math>                  (ii) Convert the following CFG to PDA by empty stack. (7)  <math>S \rightarrow 0S1/A;</math>  <math>A \rightarrow 1A0/S/\epsilon</math> <b>Infer</b> whether 0101 belongs to <math>N(M)</math>.</p>	<p>BTL-4</p>	<p>Analyze</p>
<p>4.</p>	<p>(i) If <math>L</math> is a CFL then prove that there exists PDA <math>M</math>, such that <math>L = N(M)</math>, language accepted by empty stack. (7)                  (ii) <b>Construct</b> a PDA empty store, <math>L = \{a^m b^n \mid n &lt; m\}</math>. (8)</p>	<p>BTL-6</p>	<p>Create</p>

**UNIT IV PROPERTIES OF CONTEXT FREE LANGUAGES**

Normal Forms for CFG – Pumping Lemma for CFL – Closure Properties of CFL – Turing Machines – Programming Techniques for TM.

**PART - A**

Q.No	Questions	BT Level	Competence
1.	<b>Conclude</b> the procedure for converting CNF to GNF with an example	BTL 2	Understand
2.	<b>Illustrate</b> the Basic Turing Machine model and explain in one move. What are the actions take place in TM?	BTL 3	Apply
3.	<b>Define</b> the two normal forms of CFG	BTL 1	Remember
4.	<b>Point</b> out the hierarchy summarized in the Chomsky hierarchy..	BTL 4	Analyze
5.	<b>Define</b> the pumping Lemma for CFLs	BTL1	Remember
6.	Define Turing Machine.	BTL1	Remember
7.	<b>Discuss</b> the applications of Turing machine.	BTL 2	Understand
8.	<b>Define</b> Chomskian hierarchy of language.	BTL 1	Remember
9.	What is the class of language for which the TM has both accepting and rejecting configuration? Can this be called a Context free Language? <b>Discuss.</b>	BTL 2	Understand
10.	<b>Show</b> the following grammar into an equivalent one with no unit productions and no useless symbols $S \rightarrow ABA$ $A \rightarrow aAA aBC bB$ $B \rightarrow A bB Cb$ $C \rightarrow CC cC$	BTL 3	Apply
11.	<b>Explain</b> the special features of TM? Define universal TM. Define Instantaneous description of TM	BTL 5	Evaluate
12.	<b>Define</b> GNF.	BTL 1	Remember
13.	<b>Prepare</b> the difference between finite automata and turing machine.	BTL 6	Create
14.	<b>List</b> the three ways to simplify a context free grammar. What are the properties of the CFL generated by a CFG?	BTL 5	Evaluate
15.	Draw a transition diagram for a Turing machine to <b>identify</b> $n \text{ mod } 2$ .	BTL 1	Remember
16.	<b>Express</b> the techniques for TM construction.	BTL 2	Understand
17.	<b>Develop</b> the short notes on two-way infinite tape TM.	BTL 6	Create
18.	Differentiate TM and PDA.	BTL 4	Analyze
19.	<b>Point out</b> the role of checking off symbols in a Turing Machine	BTL 4	Analyze
20.	<b>Illustrate</b> Halting Problem.	BTL 3	Apply

<b>PART - B</b>			
<b>1.</b>	<p><b>Express</b> the following grammar G into Greibach Normal Form(GNF) (13)</p> <p><math>S \rightarrow XA BB</math>  <math>B \rightarrow b SB</math>  <math>X \rightarrow b</math>  <math>A \rightarrow a</math></p>	BTL 1	Remember
<b>2.</b>	Use the CFL pumping lemma to show <b>how</b> each of these languages not to be context-free $\{ a^i b^j c^k \mid i < j < k \}$ (13)	BTL 2	Understand
<b>3.</b>	<p><b>(i)Discuss</b> a TM to accept the language <math>LE = \{ 1^n 2^n 3^n \mid n \geq 1 \}</math> (6)</p> <p><b>(ii)Construct</b> a Turing machine that <b>estimate</b> unary multiplication (Say <math>111 \times 11 = 11111</math>) (7)</p>	BTL 2	Understand
<b>4.</b>	<p><b>(i)Illustrate</b> the Turing machine for computing <math>f(m, n) = m - n</math> ( proper subtraction). (7)</p> <p><b>(ii)Demonstrate</b> a Turing Machine to compute <math>f(m+n) = m+n</math>, <math>m, n \geq 0</math> and simulate their action on the input 0100. (6)</p>	BTL 3	Apply
<b>5.</b>	<p><b>(i)Examine</b> the role of checking off symbols in a Turing Machine.(6)</p> <p><b>(ii)Describe</b> a Turing Machine M to implement the function “multiplication” using the subroutine copy(7)</p>	BTL 1	Remember
<b>6.</b>	<p><b>(i)Demonstrate</b> the implications of halting problem.(7)</p> <p><b>(ii)Show</b> that if a language is accepted by a multitape Turing machine ,it is accepted by a single-tape TM .(6)</p>	BTL 3	Apply
<b>7.</b>	<p><b>(i)Summarize</b> in detail about multihead and multitape TM with an example.(7)</p> <p><b>(ii)Construct</b> a Turing Machine to accept palindromes in an alphabet set <math>\Sigma = \{a, b\}</math>. Trace the strings “abab” and “baab”.(6)</p>	BTL 5	Evaluate
<b>8.</b>	<p><b>(i)Explain</b> the TM as computer of integer function with an example.(7)</p> <p><b>(ii)Design</b> a TM to implement the function <math>f(x) = x + 1</math>. (6)</p>	BTL 4	Analyze
<b>9.</b>	<p><b>(i)Design</b> a TM to accept the set of all strings <math>\{0, 1\}^*</math> with 010 as substring.(7)</p> <p><b>(ii)Write</b> shot notes on Two –way infinite tape TM.(6)</p>	BTL 6	Create
<b>10.</b>	<p><b>(i)Describe</b> computing a partial function with a TM.(6)</p> <p><b>(ii)Design</b> a TM to accept the language <math>LE = \{ a^n b^n c^n \mid n &gt; 1 \}</math>.(7)</p>	BTL 1	Remember
<b>11.</b>	<p><b>(i)Define</b> Turing machine for computing <math>f(m, n) = m * n, n \in \mathbb{N}</math>. (7)</p> <p><b>(ii) Write</b> notes on Partial solvability.(6)</p>	BTL -1	Remember
<b>12.</b>	<p><b>(i) Construct</b> a TM to reverse the given string <math>\{ abb \}</math>. (6)</p> <p><b>(ii) Explain</b> Multi tape and Multi head Turing machine with suitable example.(7)</p>	BTL 2	Understand

13.	(i) <b>Analyze</b> and Construct a TM to compute a function $f(w) = W^R$ where $W \in \{a,b\}$ .(7) (ii) <b>Construct</b> Turing machine (TM) that replace all occurrence of 111 by 101 from sequence of 0's and 1's.(6)	BTL 4	Analyze
14.	(i) <b>Infer</b> the Chomsky grammar classification with necessary example. (6) (ii) <b>Explain</b> a TM with no more than three states that accepts the language. $a(a+b)^*$ . Assume $\epsilon = \{a,b\}$ . (7)	BTL 4	Analyze
<b>PART – C</b>			
1.	(i) <b>Compose</b> the limitation of automata for Type 3, Type 2, type 0 languages.(7) (ii) Consider two-tape Turing Machine (TM) and determine whether the TM always writes a nonblank symbol on its second tape during the computation on any input string 'w'. <b>Formulate</b> this problem as a language and show it is undecidable.(8)	BTL-6	Create
2.	i) Define pumping lemma for CFL. Show that $L = \{a^i b^j c^k, i < j < k\}$ is not context free and <b>Judge</b> your answer.(6) ii) Construct a TM to move an input string over the alphabet $A = \{a\}$ to the right one cell. <b>Consider</b> that the tape head starts somewhere on a blank cell to the left of the input string to the right one cell, leaving all the remaining cell blank.(9)	BTL-5	Evaluate
3.	(i) Design and <b>explain</b> a TM to compute $f(m,n) = m * n$ , for all $m, n \in \mathbb{N}$ .(6) (ii) <b>Explain</b> how a multi track in a TM can be used for testing given positive integer is a prime or not(9).	BTL-4	Analyze
4.	(i) <b>Prepare</b> a subroutine to move a TM head from its current position to the right, skipping over all 0's until reaching a 1 or a blank. If the current position does not hold 0, then the TM should halt. You may assume that there are no tape symbol other than 0,1 and B(blank). Then, use this subroutine to design to TM that accepts all strings of 0's and 1's that do not have two 1's in a row.(8) (ii) <b>Write</b> short notes on checking off symbols(7)	BTL-6	Create

## UNIT V UNDECIDABILITY

Non Recursive Enumerable (RE) Language – Undecidable Problem with RE – Undecidable Problems about TM – Post's Correspondence Problem, The Class P and NP

## PART - A

Q.No	Questions	BT Level	Competence
1.	<b>Distinguish</b> between PCP and MPCP? What are the concepts used in UTMs?	BTL 2	Understand
2.	<b>List</b> out the features of universal turing machine.	BTL 1	Remember
3.	When a recursively enumerable language is said to be recursive? <b>Discuss</b> on it.	BTL 2	Understand
4.	<b>Compare</b> and contrast recursive and recursively enumerable languages	BTL 4	Analyze
5.	State <b>when</b> a problem is said to be decidable and give an example of an undecidable problem.	BTL 1	Remember
6.	<b>Define</b> NP hard and NP completeness problem.	BTL 1	Remember
7.	<b>Define</b> a universal language $L_u$ ?	BTL 1	Remember
8.	Is it true that the language accepted by a non-deterministic Turing Machine is different from recursively enumerable language? <b>Judge</b> your answer.	BTL 5	Evaluate
9.	<b>Formulate</b> the two properties of recursively enumerable sets which are undecidable	BTL 6	Create
10.	When a problem is said to be decidable? Give an example of undecidable problem. <b>Analyze</b> it.	BTL 4	Analyze
11.	What is (a) recursively enumerable languages (b) recursive sets? <b>Generalize</b> your answer.	BTL 6	Create
12.	<b>Define</b> the classes of P and NP.	BTL 1	Remember
13.	Is it true that complement of a recursive language is recursive? <b>Discuss</b> your answer	BTL 2	Understand
14.	<b>Describe</b> an example of an undecidable problem	BTL 1	Remember
15.	<b>Point out</b> the properties of recursive and recursive enumerable language.	BTL 4	Analyze
16.	<b>Illustrate</b> on Primitive Recursive Function.	BTL 3	Apply
17.	<b>Show</b> the Properties of Recursive Languages	BTL 3	Apply
18.	<b>Explain</b> about tractable problem.	BTL 5	Evaluate
19.	<b>Describe</b> post correspondence problem.	BTL 2	Understand
20.	<b>Illustrate</b> about Time and space complexity of TM.	BTL 3	Apply
<b>PART - B</b>			

1.	(i) <b>Describe</b> about the tractable and intractable problems.(7) (ii) <b>Identify</b> that “MPCP reduce to PCP”.(6)	BTL 1	Remember															
2.	(i) <b>Describe</b> about Recursive and Recursive Enumerable languages with example.(7) (ii)State and <b>describe</b> RICE theorem.(6)	BTL 1	Remember															
3.	(i) <b>Summarize</b> diagonalization language.(6) (ii) <b>Discuss</b> the significance of universal turing machine and also construct a turing machine to add two numbers and encode it .(7)	BTL 2	Understand															
4.	<b>Discuss</b> post correspondence problem .Let $\Sigma=\{0,1\}$ .Let A and B be the lists of three strings each ,defined as <table border="1" style="margin-left: 40px;"> <thead> <tr> <th></th> <th>List A</th> <th>List B</th> </tr> </thead> <tbody> <tr> <td>i</td> <td>w<sub>i</sub></td> <td>x<sub>i</sub></td> </tr> <tr> <td>1</td> <td>1</td> <td>111</td> </tr> <tr> <td>2</td> <td>10111</td> <td>10</td> </tr> <tr> <td>3</td> <td>10</td> <td>0</td> </tr> </tbody> </table> (i)Does the PCP have a solution?(7) (ii)Prove that the universal language is recursively enumerable.(6)		List A	List B	i	w <sub>i</sub>	x <sub>i</sub>	1	1	111	2	10111	10	3	10	0	BTL 2	Understand
	List A	List B																
i	w <sub>i</sub>	x <sub>i</sub>																
1	1	111																
2	10111	10																
3	10	0																
5.	(i) <b>Explain</b> computable functions with suitable example.(6) (ii) <b>Explain</b> in detail notes on Unsolvble Problems.(7)	BTL 4	Apply															
6.	(i) <b>Describe</b> in detail notes on universal Turing machines with example.(7) (ii) <b>Collect</b> and write the short notes on NP-complete problems.(6)	BTL 1	Remember															
7.	(i) <b>Show</b> that the diagonalization language ( $L_d$ ) is not a recursively enumerable.(7) (ii) <b>Illustrate</b> about unsolvability.(6)	BTL 3	Apply															
8.	(i) <b>Compare</b> the difference between recursive and recursively enumerable languages.(7) (ii) <b>Explain</b> about PCP.(6)	BTL 5	Evaluate															
9.	(i) Explain about Universal Turing machine and show that the universal language ( $L_u$ ) is recursively enumerable but not recursive. <b>Generalize</b> your answer(8) (ii) <b>Design</b> and explain how to measure and classify complexity.(5)	BTL 6	Create															
10.	(i) <b>Explain</b> about the recursively Enumerable Language with example.(6) (ii) <b>Point out</b> that the following problem is undecidable. Given two CFGs G1 and G2 is $L(G1) \cap L(G2) = \emptyset$ .(7)	BTL 4	Analyze															
11.	(i) <b>Show</b> that the characteristic function of the set of all even number is recursive .(7) (ii) <b>Illustrate</b> in detail notes on primitive recursive functions with examples.(6)	BTL-3	Apply															



12.	(i) <b>Point out</b> the Measuring and Classifying Complexity.(7) (ii)Does PCP with two lists $x=(b, b ab^3, ba)$ and $y=(b^3, ba, a)$ have a solution. <b>Analyze</b> your answer.(6)	BTL-4	Analyze
13.	(i) <b>Discuss</b> in detail about Time and Space Computing of a Turing Machine(6) (ii) <b>Express</b> two languages which are not recursively enumerable.(7)	BTL-2	Understand
14.	(i) <b>Describe</b> in detail Polynomial Time reduction and NP-completeness.(7) (ii) <b>List</b> out the short notes on NP-Hard Problems.(6)	BTL 1	Remember
<b>PART – C</b>			
1.	<b>Consider</b> and find the languages obtained from the following operations: (i)Union of two recursive languages.(5) (ii)Union of two recursively enumerable languages.(5) (iii) $L$ if $L$ and complement of $L$ are recursively enumerable.(5)	BTL-5	Evaluate
2.	Prove that the universal language is recursively enumerable but not recursive. <b>Generalize</b> your answer.(15)	BTL-6	Create
3.	(i) <b>Plan</b> and explain on decidable and un-decidable problems with an example(7) (ii) <b>Design</b> and prove that for two recursive languages $L_1$ and $L_2$ their union and intersection is recursive.(8)	BTL-6	Create
4.	(i) <b>Compare</b> and write about tractable and untractable problems with an example.(10) (ii) <b>Explain</b> the language $L_u$ and show that $L_u$ is RE language.(5)	BTL-4	Analyze